

Illuminating the I/O Optimization Path of Scientific Applications

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ISC HPC 2023



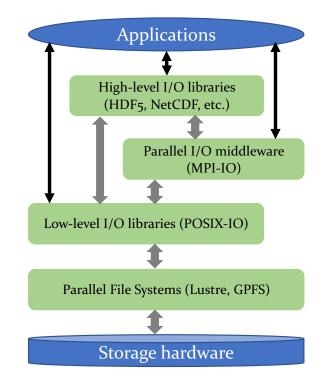


HPC I/O stack – Complex interdependencies among layers

• HPC I/O stack → complex

• Large tuning parameter space

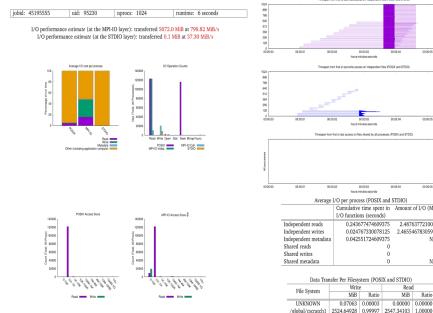
• **I/O profiling tools** for understanding I/O performance

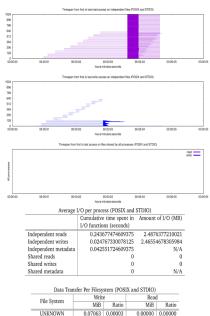


Understanding I/O performance – Darshan and DXT

- **Darshan** is a lightweight HPC I/O profiling tool
- **Darshan Extended Trace (DXT)**

- Fine grain view of the application behavior
- Interface (POSIX or MPI-IO), operation (read/write) 0
- MPI Rank, segment, offset, request size 0
- Start and end timestamp \cap
- **Challenge:** How to visualize and extract insights from DXT data?
 - Identify I/O bottlenecks Ο
 - Optimize the application 0





8_benchmark_parallel (8/6/2021)

/global/cscratch1/sd/jeanbez/paper/openPMD-api-build/bin/8.benchmark.paralle

8_benchmark_parallel (8/6/2021)

/global/cscratch1/sd/ieanbez/paper/openPMD-api-build/bin/8.benchmark.parallel

3 of 4

DXT POSIX module data

DXT, file_id: 5076057741753365924, file_name: /global/cscratch1/sd/tonglin/data_e2e/3d 28 16 16 32 32 32-36745115-1-nodes.nc4

DXT, rank: 0, hostname: nid00604

DXT, write_count: 10249, read_count: 0

DXT, mnt_pt: /global/cscratch1, fs_type: lustre

DXT, Lustre stripe_size: 16777216, Lustre stripe_count: 244

DXT, Lustre OST obdidx: 132 52 146 214 86 200 176 24 16 6 224 76 90 198 190 112 114 58 78 102 74 32 68 36 48 208 30 194 238 182 126 96
28 142 188 34 44 22 164 54 140 92 110 20 156 62 72 150 84 144 94 128 38 202 8 148 134 158 186 98 46 138 154 168 108 82 106 80 0 136 210
118 4 10 40 14 184 196 172 18 12 174 116 162 64 120 50 166 56 26 192 180 178 104 170 124 42 122 152 130 70 100 160 88 247 243 227 219 215
177 233 221 223 207 89 229 91 213 237 199 205 245 209 193 155 189 123 149 211 169 235 145 201 81 157 21 97 165 175 179 143 161 31 53 41
181 231 225 183 67 129 119 85 71 77 5 29 107 61 9 113 11 147 103 13 111 133 33 63 121 127 141 35 93 101 109 75 23 99 117 167 49 185 115 7
135 3 57 95 43 27 191 1 163 51 15 153 187 55 151 239 79 25 137 47 217 17 39 59 171 69 173 37 203 125 131 87 19 195 65 45 139 105 241 83
159 73 197 2 216 234 218 222 246 220 226 232 244 206 212 236 228 240 242

133 13 131	2 210	234 210	222 240 220	220 232 244	200 212 250 220	240 242		
# Module	Rank	Wt/Rd	Segment	0ffset	Length	Start(s)	End(s) [0	ST]
X_POSIX	0	write	0	0	1955	0.1331	0.1359	[132]
X_POSIX	0	write	1	3758106112	2038	0.1360	0.1372	[83]
X_POSIX	0	write	2	11274300928	1978	0.1372	0.1384	[7]
X_POSIX	0	write	3	18790497792	4096	0.1384	0.1391	[41]
X_POSIX	0	write	4	18790501888	366	0.1391	0.1391	[41]
X_POSIX	0	write	5	18790502254	366	0.1391	0.1392	[41]
X_POSIX	0	write	6	18790502922	150	0.1392	0.1392	[41]
X_POSIX	0	write	7	18790503072	366	0.1392	0.1392	[41]
X_POSIX	0	write	8	9728	256	0.6889	0.6894	[132]
X_POSIX	0	write	9	13824	256	0.6894	0.6922	[132]
X_POSIX	0	write	10	17920	256	0.6922	0.6926	[132]
X_POSIX	0	write	11	22016	256	0.6926	0.6930	[132]
X_POSIX	0	write	12	26112	256	0.6930	0.6937	[132]
X_POSIX	0	write	13	30208	256	0.6937	0.6942	[132]
X_POSIX	0	write	14	34304	256	0.6943	0.6946	[132]
X_POSIX	0	write	15	38400	256	0.6946	0.6951	[132]
X_POSIX	0	write	16	42496	256	0.6951	0.6956	[132]
X_POSIX	0	write	17	46592	256	0.6956	0.6961	[132]
X_POSIX	0	write	18	50688	256	0.6961	0.6966	[132]
X_POSIX	0	write	19	54784	256	0.6966	0.6970	[132]
V DOCTV	A	write	20	20000	256	A 607A	0 6074	[122]
	<pre># Module X_POSIX X_POSIX</pre>	X_POSIX 0 X_POSIX 0	# ModuleRankWt/RdX_POSIX0write	<pre># Module Rank Wt/Rd Segment X_POSIX 0 write 0 X_POSIX 0 write 1 X_POSIX 0 write 1 X_POSIX 0 write 2 X_POSIX 0 write 3 X_POSIX 0 write 4 X_POSIX 0 write 5 X_POSIX 0 write 5 X_POSIX 0 write 6 X_POSIX 0 write 7 X_POSIX 0 write 8 X_POSIX 0 write 9 X_POSIX 0 write 9 X_POSIX 0 write 10 X_POSIX 0 write 11 X_POSIX 0 write 11 X_POSIX 0 write 11 X_POSIX 0 write 13 X_POSIX 0 write 14 X_POSIX 0 write 14 X_POSIX 0 write 15 X_POSIX 0 write 16 X_POSIX 0 write 17 X_POSIX 0 write 18 X_POSIX 0 write 18 X_POSIX 0 write 19</pre>	# Module Rank Wt/Rd Segment Offset X_POSIX 0 write 0 0 X_POSIX 0 write 1 3758106112 X_POSIX 0 write 1 3758106112 X_POSIX 0 write 1 3758106112 X_POSIX 0 write 2 11274300928 X_POSIX 0 write 3 18790497792 X_POSIX 0 write 4 18790501888 X_POSIX 0 write 5 18790502254 X_POSIX 0 write 5 18790502254 X_POSIX 0 write 6 18790502254 X_POSIX 0 write 7 18790503072 X_POSIX 0 write 8 9728 X_POSIX 0 write 10 17920 X_POSIX 0 write 11 22016 X_POSIX 0	# Module Rank Wt/Rd Segment Offset Length X_POSIX 0 write 0 0 1955 X_POSIX 0 write 1 3758106112 2038 X_POSIX 0 write 2 11274300928 1978 X_POSIX 0 write 3 18790497792 4096 X_POSIX 0 write 4 18790501888 366 X_POSIX 0 write 5 18790502254 366 X_POSIX 0 write 6 18790503072 366 X_POSIX 0 write 7 18790503072 366 X_POSIX 0 write 8 9728 256 X_POSIX 0 write 9 13824 256 X_POSIX 0 write 10 17920 256 X_POSIX 0 write 12 26112 256 X_POSIX 0 <td># Module Rank Wt/Rd Segment Offset Length Start(s) X_POSIX 0 write 0 1955 0.1331 X_POSIX 0 write 1 3758106112 2038 0.1360 X_POSIX 0 write 2 11274300928 1978 0.1372 X_POSIX 0 write 3 18790497792 4096 0.1384 X_POSIX 0 write 4 18790501888 366 0.1391 X_POSIX 0 write 5 18790502254 366 0.1391 X_POSIX 0 write 6 18790503072 366 0.1392 X_POSIX 0 write 7 18790503072 366 0.1392 X_POSIX 0 write 9 13824 256 0.6889 X_POSIX 0 write 10 17920 256 0.6922 X_POSIX 0 write 11 22016 256 0.6926 X_POSIX 0 write 13<!--</td--><td># Module Rank Wt/Rd Segment Offset Length Start(s) End(s) [0] X_POSIX 0 write 0 0 1955 0.1331 0.1359 X_POSIX 0 write 1 3758106112 2038 0.1360 0.1372 X_POSIX 0 write 2 11274300928 1978 0.1372 0.1384 X_POSIX 0 write 3 18790497792 4096 0.1384 0.1391 X_POSIX 0 write 4 18790502254 366 0.1391 0.1392 X_POSIX 0 write 6 18790502252 150 0.1392 0.1392 X_POSIX 0 write 7 18790503072 366 0.1392 0.1392 X_POSIX 0 write 9 13824 256 0.6889 0.6894 X_POSIX 0 write 10 17920 256 0.6930 0.6937 X_POSIX 0 write 12 26112 256 0.6930</td></td>	# Module Rank Wt/Rd Segment Offset Length Start(s) X_POSIX 0 write 0 1955 0.1331 X_POSIX 0 write 1 3758106112 2038 0.1360 X_POSIX 0 write 2 11274300928 1978 0.1372 X_POSIX 0 write 3 18790497792 4096 0.1384 X_POSIX 0 write 4 18790501888 366 0.1391 X_POSIX 0 write 5 18790502254 366 0.1391 X_POSIX 0 write 6 18790503072 366 0.1392 X_POSIX 0 write 7 18790503072 366 0.1392 X_POSIX 0 write 9 13824 256 0.6889 X_POSIX 0 write 10 17920 256 0.6922 X_POSIX 0 write 11 22016 256 0.6926 X_POSIX 0 write 13 </td <td># Module Rank Wt/Rd Segment Offset Length Start(s) End(s) [0] X_POSIX 0 write 0 0 1955 0.1331 0.1359 X_POSIX 0 write 1 3758106112 2038 0.1360 0.1372 X_POSIX 0 write 2 11274300928 1978 0.1372 0.1384 X_POSIX 0 write 3 18790497792 4096 0.1384 0.1391 X_POSIX 0 write 4 18790502254 366 0.1391 0.1392 X_POSIX 0 write 6 18790502252 150 0.1392 0.1392 X_POSIX 0 write 7 18790503072 366 0.1392 0.1392 X_POSIX 0 write 9 13824 256 0.6889 0.6894 X_POSIX 0 write 10 17920 256 0.6930 0.6937 X_POSIX 0 write 12 26112 256 0.6930</td>	# Module Rank Wt/Rd Segment Offset Length Start(s) End(s) [0] X_POSIX 0 write 0 0 1955 0.1331 0.1359 X_POSIX 0 write 1 3758106112 2038 0.1360 0.1372 X_POSIX 0 write 2 11274300928 1978 0.1372 0.1384 X_POSIX 0 write 3 18790497792 4096 0.1384 0.1391 X_POSIX 0 write 4 18790502254 366 0.1391 0.1392 X_POSIX 0 write 6 18790502252 150 0.1392 0.1392 X_POSIX 0 write 7 18790503072 366 0.1392 0.1392 X_POSIX 0 write 9 13824 256 0.6889 0.6894 X_POSIX 0 write 10 17920 256 0.6930 0.6937 X_POSIX 0 write 12 26112 256 0.6930

DXT_MPIIO module data

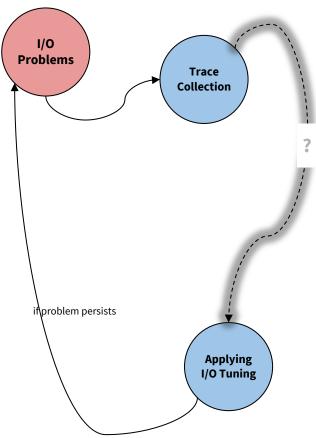
# DXT, ran # DXT, wri	nk: 0, h te_coun	ostname t: 12,	e: nid00604 read_count:	0	lobal/ <u>cscrat</u>	ch1/sd/tongl	in/data_e2e/ <u>3d_28_16_16_32_32_32-36745115-1-nodes.nc4</u>
# DAT, MMT			Segment	_type: lustre Length S	Start(s)	End(s)	
			5	5			
X_MPIIO	0	write	0	331	0.1315	0.1392	
X_MPIIO	0	write	1	262144	0.6802	1.1448	
X_MPIIO	0	write	2	262144	1.1451	1.7255	
X_MPIIO	0	write	3	262144	1.7257	2.3791	
X_MPIIO	0	write	4	262144	2.3794	3.0459	
X_MPIIO	0	write	5	262144	3.0462	4.2975	
X_MPIIO	0	write	6	262144	4.2978	5.4152	
X_MPIIO	0	write	7	262144	5.4154	6.3356	
X_MPIIO	0	write	8	262144	6.3358	7.0600	
X_MPIIO	0	write	9	262144	7.0602	7.8717	
X_MPIIO	0	write	10	262144	7.8719	8.7132	
X_MPIIO	0	write	11	96	9.5743	9.5746	

More details on Darshan Utilities:

https://www.mcs.anl.gov/research/projects/darshan/docs/darshanutil.html

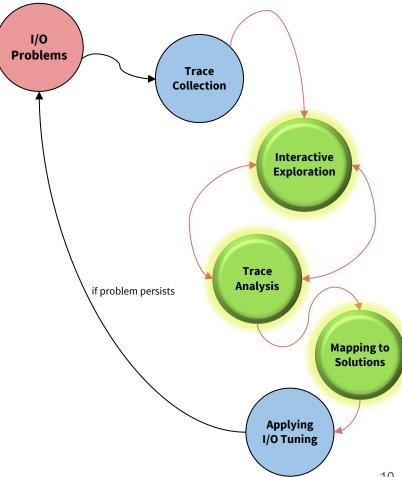
Missing dots in tuning I/O

- Despite the availability of fine grain traces
 - Gaps between trace collection, analysis, and tuning



Closing the translation gap

- A solution to close this gap requires
 - Analysis of collected metrics and traces
 - Diagnosis of root causes of poor performance
 - Recommend performance improving solutions

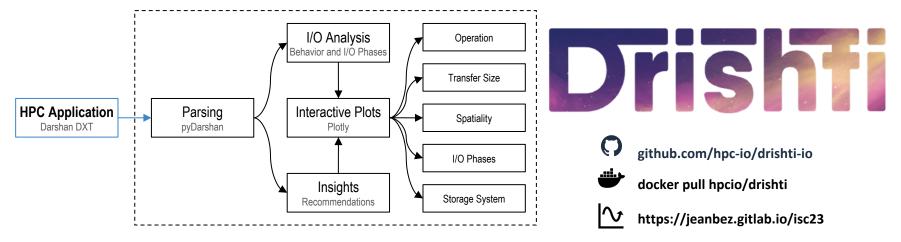


Envisioning a solution – Based on visualization and guidance

- Provide **interactive visualization** based on I/O trace files
- Display **contextual information** about I/O calls
- Understand how the application issues its **I/O requests over time**
- **Observe transformations** as the requests traverse the I/O software stack
- Detect and characterize the **distinct I/O phases** of an application
- Understand how the **ranks access the file system** in I/O operations
- Provide an extensible **community-driven framework**
- Identify and highlight common root causes of I/O performance problems
- Provide a **set of actionable items** based on the detected I/O bottlenecks

Drishti: Guiding end-users in the I/O optimization journey

- Sanskrit word; meaning 'focused gaze'
- Interactive web based log analysis framework to visualize Darshan DXT logs
 - Pinpoint root causes of I/O performance problems
 - Provide a set of actionable recommendations



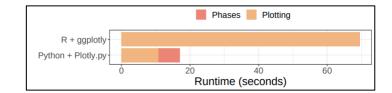
Extracting I/O Behavior from Traces

- Command line solution named darshan-dxt-parser
 - Parsed data -> Textual format -> CSV
- As an alternative, we explore the novel **PyDarshan**
 - Provides interface to binary Darshan log files -> Pandas dataframe
 - PyDarshan has shortcomings too
 - Loop over all the ranks to get a rectangular structure

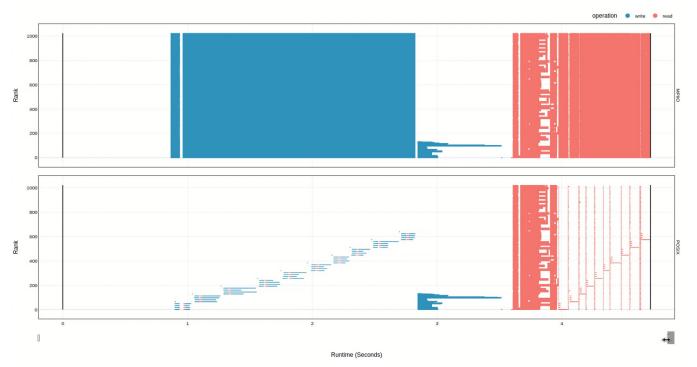


Interactively Exploring I/O Behavior

- **Static plots** limited in the information they represent
 - Space constraints
 - Pixel resolution
- We consider two solutions
 - **R** using **ggplot2**
 - **Python** using **plotly**
- Opted to rely on PyDarshan and plotly
 - Performance speedup
 - Modularity

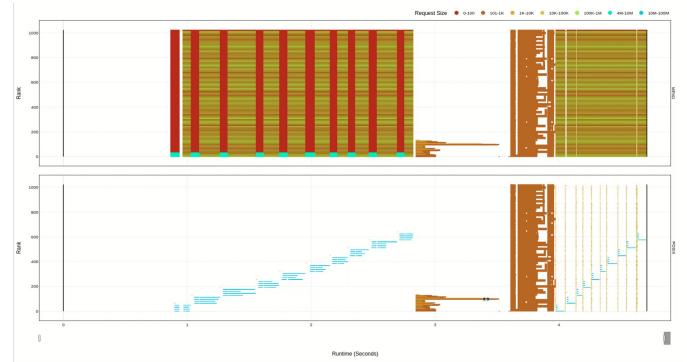


I/O Operations



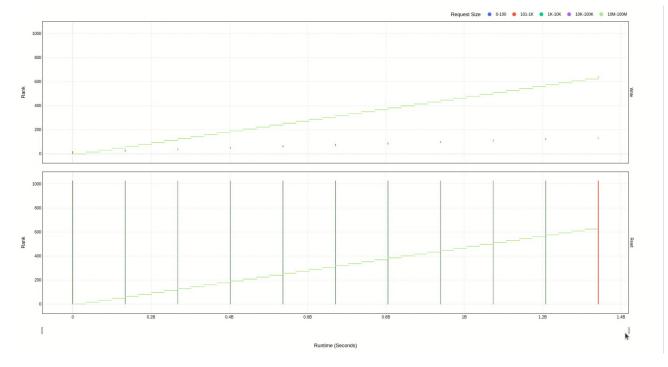
Explore the timeline by **zooming in and out** and observing how the **MPI-IO** calls are translated to the **POSIX** layer. Visualize relevant information in the context of each I/O call (rank, operation, duration, request size, and OSTs if Lustre) by hovering over a given operation.

Data Transfers



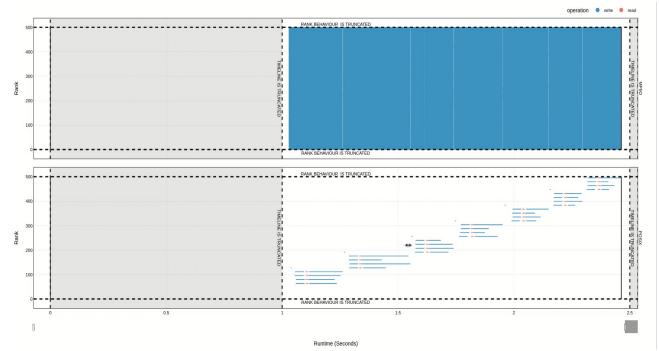
Explore the operations by size in **POSIX** and **MPI-IO**. You can, for instance, identify small or metadata operations from this visualization.

Spatiality



Explore the **spatiality** of accesses in file by each rank with contextual information.

Focused Operation View



Explore the timeline by **zooming in and out** and observing how the **MPI-IO** calls are translated to the **POSIX** layer. Visualize relevant information in the context of each I/O call (rank, operation, duration, request size, and OSTs if Lustre) by hovering over a given operation.

Automatic Detection of I/O Bottlenecks

- Variety of tools that seek to analyze performance
 - Neither provide support for auto detection of I/O bottlenecks
- Drishti provides framework which provides
 - Actionable set of recommendations in form of a report
 - Multi layered plots to diagnose and highlight bottlenecks

Root Causes	Darshan	DXT	System	Drishti
Too many I/O phases	~	~	×	~
Stragglers in each I/O phase	~	~	×	~
Bandwidth limited by a single OST I/O bandwidth	×	×	~	×
Limited by the small data size	\sim	~	×	~
Rank 0 heavy-workload	~	\checkmark	×	~
Unbalanced I/O workload among MPI ranks	~	\checkmark	×	~
Large number of small I/O requests	~	~	×	~
Unbalanced I/O workload on OSTs	~	~	~	~
Bad file system weather	×	×	~	×
Redundant/overlapping I/O accesses	~	\checkmark	×	~
I/O resource contention at OSTs	×	×	~	×
Heavy metadata load	~	×	×	~

Drishti Reports

- Relies on **counters** available in Darshan profiling logs
- Detects typical I/O performance pitfalls
 - Based on **32 checks** covering common I/O performance pitfalls
- Insights classified into **four** categories
 - **HIGH, WARN, OK, INFO**
- Provides recommendations in form of a **report**

IGH	STDIO	High STDIO usage* ($> 10\%$ of total transfer size uses STDIO)					
OK	POSIX	High number [*] of sequential read operations ($\geq 80\%$)					
OK	POSIX	High number [*] of sequential write operations ($\geq 80\%$)					
INFO	POSIX	Write operation count intensive [*] (> 10% more writes than reads)					
INFO	POSIX	Read operation count intensive [*] (> 10% more reads than writes)					
INFO	POSIX	Write size intensive* (> 10% more bytes written then read)					
INFO	POSIX	Read size intensive* (> 10% more bytes read then written)					
WARN	POSIX	Redundant reads					
WARN	POSIX	Redundant writes					
HIGH	POSIX	High number [*] of small [†] reads (> 10% of total reads)					
HIGH	POSIX	High number [*] of small [†] writes (> 10% of total writes)					
HIGH	POSIX	High number* of misaligned memory requests (> 10%)					
HIGH	POSIX	High number [*] of misaligned file requests $(> 10\%)$					
HIGH	POSIX	High number [*] of random read requests (> 20%)					
HIGH	POSIX	High number [*] of random write requests (> 20%)					
HIGH	POSIX	High number [*] of small [†] reads to shared-files (> 10% of reads)					
HIGH	POSIX	High number [*] of small [†] writes to shared-files (> 10% of writes)					
HIGH	POSIX	High metadata time [*] (one or more ranks spend > 30 seconds)					
HIGH	POSIX	Rank o heavy workload					
HIGH	POSIX	Data transfer imbalance between ranks (> 15% difference)					
HIGH	POSIX	Stragglers detected among the MPI ranks					
HIGH	POSIX	Time imbalance * between ranks (> 15% difference)					
WARN	MPI-IO	No MPI-IO calls detected from Darshan logs					
HIGH	MPI-IO	Detected MPI-IO but no collective read operation					
HIGH	MPI-IO	Detected MPI-IO but no collective write operation					
WARN		Detected MPI-IO but no non-blocking read operations					
WARN		Detected MPI-IO but no non-blocking write operations					
OK	MPI-IO						
OK	MPI-IO	Detected MPI-IO and collective write operations					
HIGH		Detected MPI-IO and inter-node aggregators					
WARN		Detected MPI-IO and intra-node aggregators					
OK	MPI-IO	Detected MPI-IO and one aggregator per node					

	Drishti					
DRISHTI v.0.3 -						
JOB	1190243					
EXECUTABLE	bin/8_benchmark_parallel					
DARSHAN	jlbez_8_benchmark_parallel_id1190243_7-23-45631-11755726114084236527_1.darsha					
	2021-07-23 16:40:31+00:00 to 2021-07-23 16:40:32+00:00 (0.00 hours)					
FILES	6 files (1 use STDIO, 2 use POSIX, 1 use MPI-IO)					
PROCESSES	64					
HINTS	romio_no_indep_rw=true cb_nodes=4					
1 critical issu	es, 5 warnings, and 5 recommendations					
METADATA						
METADATA						
Application i	s read operation intensive (6.34% writes vs. 93.66% reads)					
	ight have redundant read traffic (more data read than the highest offset)					
	ight have redundant read charine (more data read than the highest offset)					
► Application m	ight have redundant write trainic (more data written than the highest offset)					
OPERATIONS						
OFERATIONS						
Application i	ssues a high number (285) of small read requests (i.e., < 1MB) which represent					
	ad/write requests					
) small read requests are to "benchmark.h5"					
↔ Recommendat	lons:					
	buffering read operations into larger more contiguous ones					
⇔ Since the	appplication already uses MPI-IO, consider using collective I/O calls (e.g.					
↔ Since the MPI_File_read_al	appplication already uses MPI-IO, consider using collective I/O calls (e.g. l() or MPI_File_read_at_all()) to aggregate requests into larger ones					
Since the MPI_File_read_al ► Application m	appplication already uses MPI-IO, consider using collective I/O calls (e.g. l() or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests					
Since the MPI_File_read_al ► Application m ► Application m	appplication already uses MPI-IO, consider using collective I/O calls (e.g. () or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests					
Since the MPI_File_read_al ► Application m ► Application m	appplication already uses MPI-IO, consider using collective I/O calls (e.g. l() or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests					
Since the MPI_File_read_al ► Application m ► Application m ► Application u	appplication already uses MPI-IO, consider using collective I/O calls (e.g. [] or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests ses MPI-IO and read data using 640 (83.55%) collective operations					
Since the MPI_File_read_al ► Application m ► Application m ► Application u ► Application u	appplication already uses MPI-IO, consider using collective I/O calls (e.g. [() or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests ses MPI-IO and read data using 640 (83.55%) collective operations ses MPI-IO and write data using 768 (100.00%) collective operations					
Since the MPI_File_read_al ► Application m ► Application m ► Application u ► Application u	appplication already uses MPI-IO, consider using collective I/O calls (e.g. () or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests ses MPI-IO and read data using 640 (83.55%) collective operations ses MPI-IO and write data using 768 (100.00%) collective operations ould benefit from non-blocking (asynchronous) reads					
<pre></pre>	apppliation already uses MPI-IO, consider using collective I/O calls (e.g. () or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) read requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests ses MPI-IO and read data using 640 (83.55%) collective operations ses MPI-IO and write data using 768 (100.00%) collective operations ould benefit from non-blocking (asynchronous) reads ions:					
 Since the MPI_File_read_al Application m Application u Application u Application u Application c Gecommendat Since you 	appplication already uses MPI-IO, consider using collective I/O calls (e.g. [() or MPI_File_read_at_all()) to aggregate requests into larger ones ostly uses consecutive (2.73%) and sequential (90.62%) write requests ostly uses consecutive (19.23%) and sequential (76.92%) write requests ses MPI-IO and write data using 640 (83.55%) collective operations ses MPI-IO and write data using 768 (100.00%) collective operations ould benefit from non-blocking (asynchronous) reads ions: use MPI-IO, consider non-blocking/asynchronous I/O operations (e.g.,					
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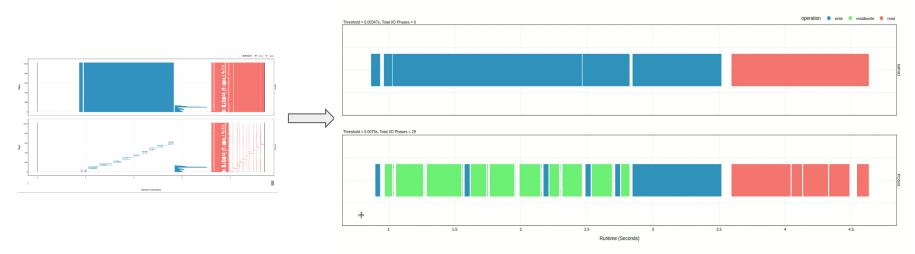
Exploring I/O Phases and Bottlenecks

- I/O phase is a **time period** where an application is accessing its data
- Factors outside an application's scope could cause an I/O phase to take longer
 - Network Interference
 - Storage system congestion
- Drishti adds an **interactive visualization** to detect I/O phases based on DXT trace data
 - Multi layered plots to detect workload imbalance and rank 0 heavy workload
- Gives a detailed picture of I/O phases and I/O patterns in the data
 - Helpful in extracting information related **bottlenecks** such as stragglers

Exploring I/O Phases and Bottlenecks (contd ...)

- Finding I/O phases is **not a trivial task** due the sheer amount of data
 - \circ \quad Millions of operations in order of milliseconds
- We use **PyRanges** to find similar and overlapping behavior between ranks
 - PyRanges is a genomics library used for handling genomic intervals
 - Use a threshold value to merge I/O phases closer to each other
- While computing I/O phases, we keep track of the **duration** between each
 I/O phase
 - **Threshold** to merge I/O phases close to each other

Exploring I/O Phases and Bottlenecks (contd ...)

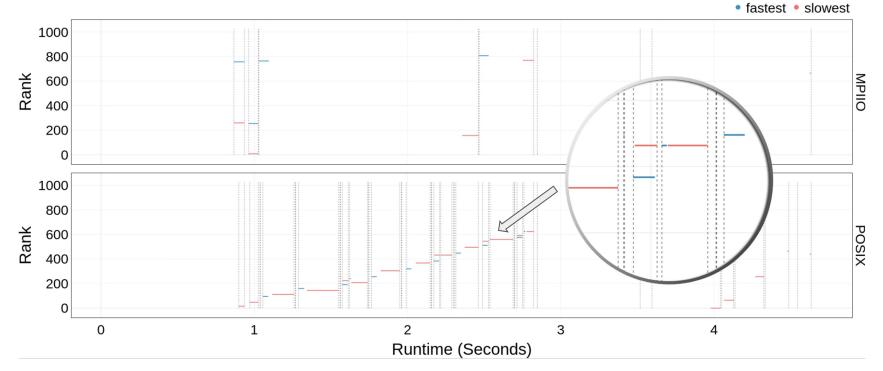


Explore the **I/O phases** of the application. Contextual information like the **fastest rank, fastest rank duration, slowest rank**, and **slowest rank duration** are available when hovering over an I/O phase.

Exploring I/O Phases and Bottlenecks (contd ...)

- I/O stragglers in each phase could define a **critical path** impairing performance
- **Exclusive plot** to highlight the I/O phases and fast and straggler in each phase
 - Handle each interface separately
- Detect **slow ranks** across the entire execution or storage servers

Stragglers



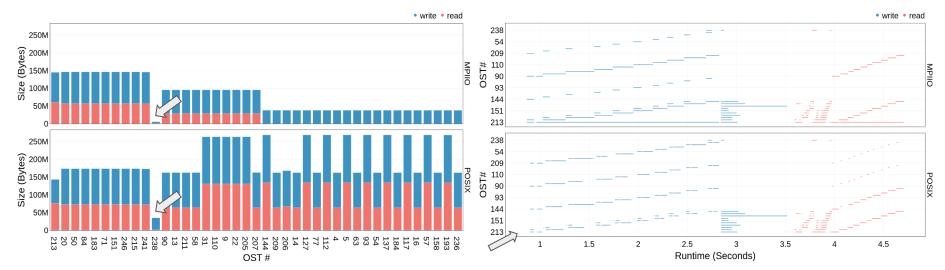
Explore the **stragglers** in the entire execution. Upon hovering over a phase, all the information related to the **fastest** and **slowest** rank is shown. The dotted lines represent the start and the end of a phase

File System Usage

- DXT captures information related to the **storage servers**
- Drishti provides a visualization to explore the OST usage of the I/O requests
- Also provides a visualization to depict data transfer sizes for each OST

Towards Exploring File System Usage (contd ...)

- The plots show
 - Very small data transfer for OST# 238 for MPI IO
 - Size increases at POSIX due to transformations as the request goes through the stack
 - OST# 213 is **most used** for MPI IO
 - OST# 238 is the least used across both MPI IO and POSIX



Explore the **file system usage** for the entire execution. The plot on the **left** shows **OST usage data transfer** sizes and the plot on the **int** shows **OST usage of I/O requests** over time

Putting Drishti into practice

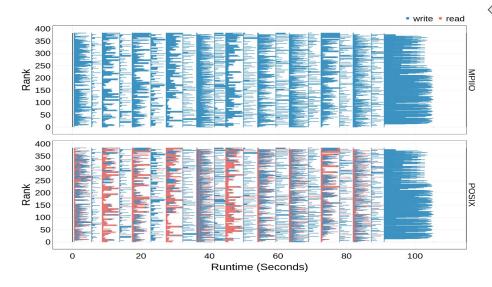
- Demonstrate Drishti to identify I/O performance bottlenecks
- Experiments conducted at:
 - Cori at NERSC
 - Summit at OLCF
- **Four** use cases:
 - OpenPMD
 - AMReX
 - E2E (available in companion repo)
 - H5Bench Write (available in companion repo)
- Used **h5bench** to generate the benchmarks

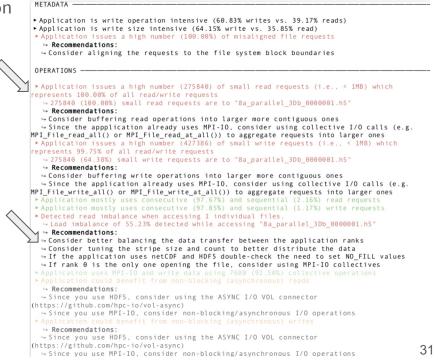
I/O Bottlenecks in OpenPMD

- Open Standard for Particle Mesh Data Files (OpenPMD)
 - Particle and mesh data in scientific simulations and experiments
- **Summit** with 64 compute nodes, 6 ranks per node, and a total of 384 MPI ranks
 - Mesh size is $[65536 \times 256 \times 256]$, 10 iterations
- For this scenario, OpenPMD takes an average **110.6 seconds**

I/O Bottlenecks in OpenPMD

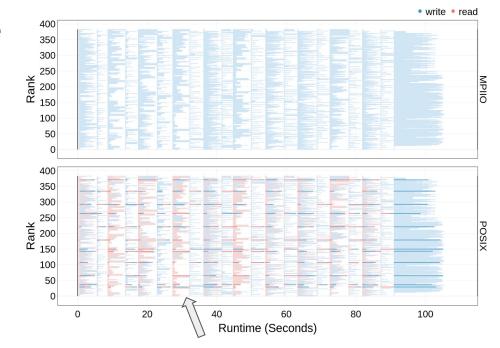
- Majority of the read and write requests are small
 - I/O calls are not using the **MPI-IO's collective** option





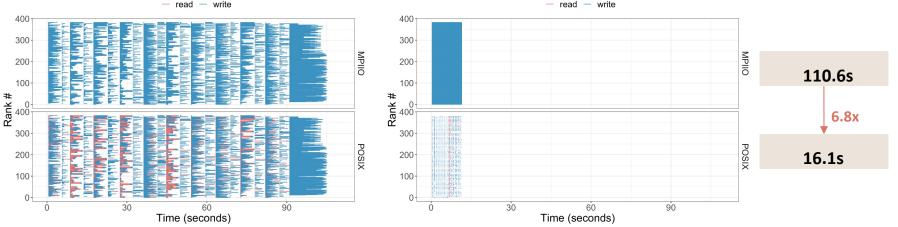
I/O Bottlenecks in OpenPMD

• Moreover, Drishti detected an **imbalance** when accessing the data



I/O Bottlenecks in OpenPMD - optimized

- Collective HDF5 metadata were not actually collective due to an issue introduced in HDF5 1.10.5
 - Fixed that issue by using **HDF5 1.10.4** (or using **1.10.6** or later) and enabling collective metadata I/O
- Drishti suggested larger buffer sizes
 - Used ROMIO hints to set the aggregators to **1 agg/node** and set the **cb_buffer_size** to 16 MB
 - Used GPFS large block I/O
- With HDF5 1.10.4 combined with other optimizations gives a total of **6.8x** speedup from baseline

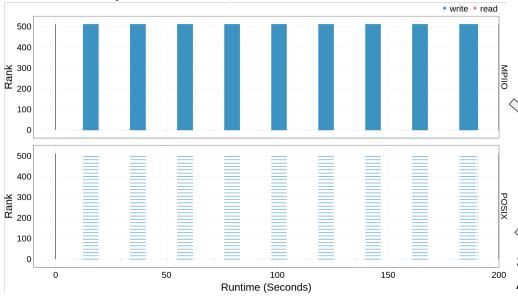


Improving AMReX with Asynchronous I/O

- AMReX is an adaptive mesh refinement (AMR) framework
 - Solves partial differential equations on block-structured meshes
 - Used by several applications in the Exascale Computing Project (ECP)
- I/O benchmark on **Cori** with 32 compute nodes, 512 ranks
 - 1024 domain size, 10 plot files, 10 seconds sleep time between writes

Improving AMReX with Asynchronous I/O (contd ...)

- The report suggests using
 - Larger buffer sizes
 - Asynchronous I/O VOL connector





17 files (3 use STDIO, 2 use POSIX, 10 use MPI-IO)

romio no inden rwatrue ch nodesad

FILES:

HTNTC .

COMPUTE NODES PROCESSES

Application is w	rite operation intensive (99	.98% writes vs. 0.02% reads)			
 Application is w 	rite size intensive (100.00%	write vs. 0.00% read)			
OPERATIONS					
		small write requests (i.e.,	< 1MB) which represents	. 99.99% of all write reque	ests
	small write requests are to				
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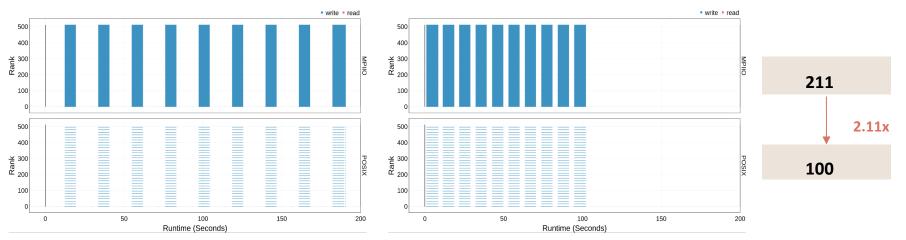
Consider buffering write operations into larger more contiguous ones



Since you use HDF5, consider using the ASYNC I/O VOL connector (https://github.com/hpc-io/vol-async)

Improving AMReX with Asynchronous I/O (contd ...)

- Added asynchronous I/O VOL Connector
 - Makes the operations **non blocking**
 - Hide time spent in time I/O while the application continues its computation
- Majority of request sizes are very small (< 1MB) for all the 10 plot files
 - Set **stripe size** to 16MB



Future Work

- Integrate additional metrics and system logs to broaden the spectrum of I/O issues
 - Global API to consume metrics from distinct sources e.g. Recorder
- Map performance optimization recommendations to the **exact source code** line numbers
 - Static code analysis
 - Modified code instead of generic snippets in Drishti reports
- Community guidelines on how to contribute to this tool
 - Aid in keeping up with the latest advancements in I/O libraries and systems
 - \circ Reach out to novel systems for support

Conclusion

- Pinpointing root causes of I/O inefficiencies requires:
 - Detailed metric **analysis**
 - **Understanding** of the HPC I/O stack
- Drishti is an interactive web based **analysis framework** which
 - Seeks to **close the gap** between trace collection, analysis, and tuning
 - **Automatically** detects common root causes of I/O performance inefficiencies
 - Provides actionable **recommendations** to the user
- Applicability demonstrated with **optimization** of OpenPMD and AMReX applications
- Companion Repository: https://jeanbez.gitlab.io/isc23





Illuminating the I/O Optimization Path of Scientific Applications

Hammad Ather, Jean Luca Bez, Boyana Norris, Suren Byna

